

Information Theory and Networks

Lecture 10: Sampling with Fair Coins

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Part I

Sampling with Fair Coins

USA Today has come out with a new survey: Apparently three out of four people make up 75 percent of the population.

David Letterman

Problem

Imagine you have a fair coin, but you want to sample from an arbitrary distribution, how would you do it?

From [CT91, p.110-116]

Example 1

Example: use a sequence of fair coin tosses to generate a random variable X with PMF

$$X = \begin{cases} a, & \text{with probability } 1/2, \\ b, & \text{with probability } 1/4, \\ c, & \text{with probability } 1/4, \end{cases}$$

Example 1

Example 1

Example 1

General Problem

- We want to generate a random variable $X \in \Omega = \{1, 2, \dots, m\}$
 - ▶ X has PMF $\{p_1, p_2, \dots, p_m\}$
- We have a series of (independent) fair coin tosses Z_1, Z_2, \dots
 - ▶ let T denote the number of coin tosses (which is potentially a RV)
 - ▶ we'd like methods that minimise $E[T]$

General Problem

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Example 1

Example 1

Problem

How could we go about designing such a tree for a *dyadic* distribution (one whose probabilities are powers of two)?

Is it related to entropy?

Example 2

Example 2

Problem

What about non-dyadic probabilities?

Obviously, we could do powers of three with ternary codes, and so on, do lets assume that the probabilities don't all fit some simple power-law.

Theorem

The expected number of fair bits $E[T]$ required by the optimal algorithm to generate a random variable X satisfies

$$H(X) \leq E[T] < H(X) + 2$$

Proof: see [CT91, pp.115-116]

Example 3

Example 3

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Example 3

Problem

What if you don't even know if your coin is fair?

We could have also looked into Stochastic Computing here – i.e., techniques for doing computation using operations stochastic processes.

Source Coding and 20 Questions

Yet another way to think about coding

- 20 questions:
 - ▶ Want to guess a 'fact' — say an experiment's outcome
 - ▶ Only allowed Yes/No questions
 - ▶ Want to find the most efficient set of questions
- Obviously, Huffman code is optimal way of generating questions if we know the PMF

Further reading I



Thomas M. Cover and Joy A. Thomas, *Elements of information theory*, John Wiley and Sons, 1991.